

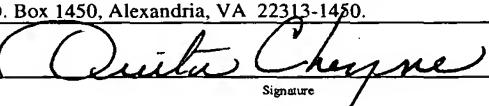
Application for United States Letters Patent

For

**ALLOCATION OF POWER AND CHANNELIZATION CODES FOR
DATA TRANSFERS**

By

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ALLOCATION OF POWER AND CHANNELIZATION CODES FOR DATA TRANSFERS

BACKGROUND OF THE INVENTION

5 1. FIELD OF THE INVENTION

This invention relates generally to the field of telecommunications and, more particularly, to wireless communications.

2. DESCRIPTION OF THE RELATED ART

Owners and/or operators of communication networks, *i.e.*, the service providers, are 10 constantly searching for methods and equipment that can meet the changing needs of their subscribers. Subscribers of communication networks, including wireless communication networks, require higher information throughput in order to exploit the expanding range of services being provided by current communication networks. For example, wireless communication subscribers are now able to have simultaneous access to data networks such 15 as the Internet and to telephony networks such as the Public Switched Telephone Network (PSTN).

Speech transmission is the main service supported by the first and second-generation wireless communication systems. However, the demand for data services is ever-increasing. Data traffic, unlike voice traffic, tends to occur in bursts and is more delay tolerant. The 20 amount of data that can be sent to a user is dependent on the amount of data that has arrived for transfer to the user. One technique for communicating data to users includes identifying the users with data packets ready for transmission, prioritizing the users in order by channel condition, and sending the data using the available channel resources.

Downlink traffic for different traffic rates is conveyed using a multiple rate combination of the basic rate. Traffic with rate m times the basic rate may be sent using m codes and the received signal power may be m times larger than the basic power allocated to the basic rates. However, in the case of a shared data and voice application, the total 5 available power at the transmitter for data traffic varies due to various reasons such as overhead power allocation or dynamic changes in the voice user's power consumption. These variations could be very slow (*e.g.*, in the order of hours) or could be fast varying (*e.g.*, as function of the channel variation in the order of milliseconds). In addition, the total system available power at the transmitter may also change over time. Hence, an assumption of fixed 10 power allocation for data service, for example, is not valid in these situations.

In Universal Mobile Telecommunications Systems (UMTS), users are code multiplexed using different orthogonal variable spreading factor (OVSF) codes and variable processing gain. Typically, the number of codes assigned to the users depends on the size of the data packet to be transferred, and the available transmit power is evenly divided amongst 15 the total number of codes. Assigning the channelization codes and then dividing the available power amongst the codes does not optimize the capacity of the data channels. In the case where the total allocated codes for a user are assigned based on the user packet or buffer size, the total power assigned is thus proportional to the total number of assigned codes rather than optimized to the channel conditions. However, users with better channel 20 conditions may not require the same power to meet a target error rate than a user with a more degraded channel.

The present invention is directed to overcoming, or at least reducing the effects of, one or more of the problems set forth above.

SUMMARY OF THE INVENTION

One aspect of the present invention is seen in a method of communication. The method includes assigning at least one channelization code to a data packet. A portion of power available for communicating is assigned to the channelization code based on a channel quality metric.

Another aspect of the present invention is seen in a system including a transmitter, and an allocation unit. The transmitter is adapted to communicate data packets. The allocation unit is adapted to assign at least one channelization code to a data packet and assign a portion of power available for communicating to the channelization code based on a channel quality metric associated with the data packet.

BRIEF DESCRIPTION OF THE DRAWINGS

The invention may be understood by reference to the following description taken in conjunction with the accompanying drawings, in which like reference numerals identify like elements, and in which:

15 Figure 1 is a simplified block diagram of a communications system in accordance with one illustrative embodiment of the present invention;

Figure 2 is a simplified flow diagram of a method of communicating in accordance with another illustrative embodiment of the present invention;

20 Figure 3 is a simplified flow diagram of one particular embodiment of a method for allocating power and channelization codes to implement the communicating of Figure 2;

Figure 4 is a simplified flow diagram of a second particular embodiment of a method for allocating power and channelization codes to implement the communicating of Figure 2; and

Figures 5A through 5C are simplified flow diagrams of methods for prioritizing users

5 and assigning resources based on different quality of service scenarios.

While the invention is susceptible to various modifications and alternative forms, specific embodiments thereof have been shown by way of example in the drawings and are herein described in detail. It should be understood, however, that the description herein of specific embodiments is not intended to limit the invention to the particular forms disclosed, 10 but on the contrary, the intention is to cover all modifications, equivalents, and alternatives falling within the spirit and scope of the invention as defined by the appended claims. Moreover, it should be emphasized that the drawings of the instant application are not to scale but are merely schematic representations, and thus are not intended to portray the specific dimensions of the invention, which may be determined by skilled artisans through 15 examination of the disclosure herein.

DETAILED DESCRIPTION OF SPECIFIC EMBODIMENTS

Illustrative embodiments of the invention are described below. In the interest of clarity, not all features of an actual implementation are described in this specification. It will of course be appreciated that in the development of any such actual embodiment, numerous 20 implementation-specific decisions must be made to achieve the developers' specific goals, such as compliance with system-related and business-related constraints, which will vary from one implementation to another. Moreover, it will be appreciated that such a

development effort might be complex and time-consuming, but would nevertheless be a routine undertaking for those of ordinary skill in the art having the benefit of this disclosure.

Referring now to Figure 1, a simplified block diagram of a communications system 100 in accordance with one embodiment of the present invention is provided. A transmitter 110 communicates with a plurality of users 120 within a sector 130 of a communications cell 140. The transmitter 110 includes a data queue 112 and an allocation unit 114. As will be described in greater detail below, the allocation unit 114 employs an iterative technique to assign channelization codes and power levels to optimize the capacity of a data channel used to communicate data with the users 120. In the illustrated embodiment, the transmitter 110 is a conventional multicode spread spectrum transmitter, the construct and operation of which are well known to those of ordinary skill in the art.

In the illustrated embodiment, the communications system 100 is a multi-cell and multi-code packet data code division multiple access (CDMA) system with three-sector cells, however, the application of the present invention is not limited to any particular system architecture. For packet data, downlink traffic is communicated over a shared channel in which each user 120 listens for indication that data currently sent in the shared channel is destined for the particular user 120. Channel quality information for the receiver (not shown) of the user 120 is fed back to the transmitter 110 through a feedback channel in the uplink. Each sector 130 has an associated transmitter 110 with an allocation unit 114.

For its sector 130, the allocation unit 114 can allocate at most N_c orthogonal channelization codes with the same spreading factor to a given user. A maximum of N_u users can receive packets over the same transmission time interval (TTI). For ease of illustration and to avoid obscuring the instant invention, it is assumed that the packet

transmission time interval is a constant. For example, the current high speed downlink packet access (HSDPA) channel implemented in the publicly available Universal Mobile Telecommunications System (UMTS) standard has a fixed spreading factor of 16 among which at most 15 channelization codes can be used to transmit one packet over a 2ms TTI.

5 HSDPA also allows more than one of the users to be code multiplexed during a TTI (*i.e.*, a user may be assigned more than one channelization code. Hybrid-ARQ (Automatic Repeat Request) operation with asynchronous incremental redundancy (IR) is performed for each transmission. The allocation unit 114 maintains a list of the active data users in the sector 130. In general, the allocation unit 114 associated with each sector 130 selects a user 120

10 based on the user data backlog, the QoS (quality of service) requirements for each packet, the available system resources, and the feedback from each user 120 listed in its user active set.

The available system resource includes the number of channelization codes left after taking into account all dedicated and common channels other than the downlink shared packet channel, such as the common pilot channel, dedicated control channel, signaling channels, *etc.*; and the available power for transmitting packet data at the time of scheduling. The user feedback includes channel quality and acknowledgement/non-acknowledgement (ACK/NACK) of the previously received packet.

The allocation unit 114 performs two major functions, first, it prioritizes the packets based on buffer backlog and QoS requirements; then it determines how many packets can be 20 served from each priority level based on the channel quality and the available system resources. The allocation unit 114 attempts to optimize system capacity in accordance with the method depicted in Figure 2, which shows a simplified flow diagram of a method for communicating in accordance with another illustrative embodiment of the present invention.

In block 200, at least one channelization code is assigned to a data packet. In block 210, a portion of power available for communicating is assigned to the channelization code based on a channel quality metric. As used herein, the term “data packet” is intended to include any set of data elements to be communicated and is not limited to any particular type of data (e.g.,

5 mathematical data, voice data, video data, raw binary data, *etc.*) or any particular protocol.

The operation of the allocation unit 114 in implementing the method of Figure 2 is now described in greater detail. The allocation unit 114 prioritizes the packets into N_p different priority classes. Assume that for each priority class, there are M_p packets waiting to be transmitted. Therefore, in a given sector 130, there are $N_p * M_p$ packets in the data queue

10 112. Note that any given user 120 can have more than one packet in the queue and each user 120 can be scheduled to receive more than one packet during the same TTI. It may be efficient to allow one of the users 120 to have more than one packet assigned under various conditions. For example, the user 120 may have a very good channel condition, and after taking out the required number of channelization codes for the first data packet transmission,

15 the same user 120 may be the best remaining user 120 among all other users 120. Hence, the second data packet to the same user 120 is scheduled using a portion of the remaining channelization codes. Another possible situation is when there are packets for the same user 120 from different priority classes. Note that if two packets from the same priority class are scheduled to transmit to the same user during the same TTI, both packets will have the same

20 power for all the channelization codes allocated to the packets, since all orthogonal channelization codes suffer the same fading and path loss for a given sector 130/user 120 pair. In some embodiments, if two packets are from different priority classes, then the two packets may be transmitted using different power levels, due to the fact that the available power for data packets from different priority classes may be different.

The allocation unit 114 determines how much power should be allocated to each channelization code and the number of channelization codes assigned to each data packet for its transmission. For the following illustration, P_{data} is the total available power for packet data for a given sector 130 and N_c represents the number of remaining channelization codes.

5 Since the allocation unit 114 strives to fulfill the QoS requirements, it first attempts to serve the packets among the highest priority class. In determining the power and channelization code assignments, the allocation unit 114 attempts to optimize the sum of the Shannon capacity for all channels in bit/s/Hz, which is defined by:

$$\frac{C}{B} = \sum_i \log_2 (1 + SNR_i), \quad (1)$$

10 where B is the system bandwidth, and the summation is over all the channelization codes used for packet data. For ease of illustration and to avoid obscuring the instant invention, it is assumed that the channels are flat Rayleigh fading channels that are quasi-static over the period of interest. The signal-to-noise ratio per code is defined by:

$$SNR_i \equiv \frac{S_i}{N_i} = \frac{|\alpha_i|^2 \beta_i P_{data}}{I_{oc} + N_o}, \quad (2)$$

15 where α_i is the complex-valued flat Rayleigh fading channel gain, β_i is the power fraction allocated to the i -th channelization code, and I_{oc} is the total interference from other cells. For purposes of this illustration, it is assumed that the QoS class for all packets is the same, and each user 120 has, at most, one packet in that QoS class. Further illustrations will address multiple QoS classes and multiple packets. The allocation unit 114 attempts to define
20 values for β_i such that Equation (1) is maximized and

$$\sum_i \beta_i = 1, \quad \beta_i \geq 0, \forall i. \quad (3)$$

Let M be the number of data users in a sector 130 and W_M be the set of channelization codes assigned to the m -th user's packet. If ϕ_m is the size of the set W_m , then Equation (3) can be written as:

$$5 \quad \sum_{m=1}^M \phi_m \chi_m = 1, \quad \chi_m \geq 0, \forall m, \quad (4)$$

where $\chi_m = \beta_i, \forall i \in W_m$ is the fraction of data power that is assigned to each of the channelization codes for user m . Note that ϕ_m is subject to the following constraint:

$$\sum_{m=1}^M \phi_m = N_c, \quad \phi_m \in \mathbb{N}. \quad (5)$$

The following cost function for the optimization problem may be defined using

10 Lagrange multipliers:

$$J(\chi_m, \phi_m) = \sum_{m=1}^M \sum_{i \in C_m} \log_2 \left(1 + \frac{|\eta_m|^2 \chi_m P_{data}}{I_{oc} + N_o} \right) + \lambda \left(\sum_{m=1}^M \phi_m \chi_m - 1 \right) + \mu \left(\sum_{m=1}^M \phi_m - N_c \right), \quad (6)$$

where $\eta_m = \alpha_i, \forall i \in C_m$ is the complex flat fading for the m -th data user. The first term of the Equation (6) can be further simplified, resulting in:

$$J(\chi_m, \phi_m) = \sum_{m=1}^M \phi_m \log_2 \left(1 + \frac{|\eta_m|^2 \chi_m P_{data}}{I_{oc} + N_o} \right) + \lambda \left(\sum_{m=1}^M \phi_m \chi_m - 1 \right) + \mu \left(\sum_{m=1}^M \phi_m - N_c \right). \quad (7)$$

15 Taking the derivative of Equation (7) with respect to χ_m , yields the following set of

equations:

$$\frac{\partial J(\chi_m, \phi_m)}{\partial \chi_m} = \frac{\phi_m \gamma_{DPR} \hat{C}_{pilot,m}}{\left(1 + \chi_m \gamma_{DPR} \hat{C}_{pilot,m}\right)} + \lambda \phi_m = 0, \quad (8)$$

where $\hat{C}_{pilot,m}$ is the channel quality estimate from the m -th user based on the measurement from the common pilots, and is defined as:

$$5 \quad \hat{C}_{pilot,m} = \frac{|\eta_m|^2 P_{pilot}}{I_{oc} + N_o}, \quad (9)$$

and γ_{DPR} is the data-to-pilot power ratio which can be expressed as:

$$\gamma_{DPR} = \frac{P_{data}}{P_{pilot}}. \quad (10)$$

The data-to-pilot power ratio is assumed to be known at the sector 130. The channel quality estimates are fed back from each user 120 to the transmitter 110 on a periodic basis.

10 Note that the channel quality estimates used in Equation (9) are the same for all codes used for the same user. Therefore, even in a multicode scenario, only one channel quality value is fed back from a given user 120. Equation (8) can be written as:

$$\frac{1}{\gamma_{DPR} \hat{C}_{pilot,m}} + \chi_m = -\frac{1}{\lambda}. \quad (11)$$

Combining Equations (4) and (11) yields:

$$15 \quad \sum_{m=1}^M \phi_m \chi_m = 1 = -\sum_{m=1}^M \frac{\phi_m}{\lambda} - \sum_{m=1}^M \frac{\phi_m}{\gamma_{DPR} \hat{C}_{pilot,m}}. \quad (12)$$

Solving Equation (12) for λ , and incorporating Equation (5) yields:

$$\lambda = \frac{-N_c}{1 + \sum_{m=1}^M \frac{\phi_m}{\gamma_{DPR} \hat{C}_{pilot,m}}}, \quad (13)$$

where λ represents an intermediate optimization parameter associated with the power constraint. Substituting Equation (13) into Equation (7) and optimizing over ϕ_m yields:

$$5 \quad \frac{\partial J(\chi_m, \phi_m)}{\partial \phi_m} = \log_2 \left(1 + \frac{|\eta_m|^2 \chi_m P_{data}}{I_{oc} + N_o} \right) - \frac{N_c \chi_m}{1 + \sum_{m=1}^M \frac{\phi_m}{\gamma_{DPR} \hat{C}_{pilot,m}}} + \mu = 0. \quad (14)$$

Multiplying both sides of Equation (14) by ϕ_m , summing over all m , and combining with Equations (4) and (5) yields:

$$\mu = \frac{1}{1 + \sum_{m=1}^M \frac{\phi_m}{\gamma_{DPR} \hat{C}_{pilot,m}}} - \frac{1}{N_c} \sum_{m=1}^M \phi_m \log_2 \left(1 + \frac{|\eta_m|^2 \chi_m P_{data}}{I_{oc} + N_o} \right), \quad (15)$$

or from Equation (13):

$$10 \quad \mu = -\frac{\lambda}{N_c} - \frac{1}{N_c} \sum_{m=1}^M \phi_m \log_2 \left(1 + \frac{|\eta_m|^2 \chi_m P_{data}}{I_{oc} + N_o} \right), \quad (16)$$

where μ represents another intermediate optimization parameter associated with the number of channelization codes.

As evidenced by Equations (14) and (16) a non-linear optimization problem exists, and a closed-form solution to this problem is difficult to find. However, for $M \leq 3$, it is

possible to solve for χ_m and ϕ_m iteratively as illustrated by the simplified flow diagram of Figure 3.

In block 300, the set of channelization codes for the users 120, $\{\phi_m^0\}$, is initialized.

The values for $\{\phi_m^0\}$ may be set at arbitrary values (e.g., even number of channelization codes 5 for each user 120 with an associated data packet). In block 310, the value of λ is determined in accordance with Equation (13). In block 320, values for the set of power fractions, $\{\chi_m^n\}$, are determined for the users 120 with assigned channelization codes in accordance with Equation (11). In block 330, the value of μ is determined in accordance with Equation (14).

In block 340, the values for the set of channelization codes, $\{\phi_m^n\}$, $\forall \phi_m^n \in \mathbb{N}$, are determined 10 in accordance with Equations (4), (5), and (16). In block 350, it is determined if the value of $\phi_m^n = \phi_m^{n+1}$, $\forall n$, where n is sufficiently large (i.e., the solution has converged). If the solution converges, the iteration terminates at block 360. Otherwise, the iterative process repeats at block 310.

The previous illustration assumes that data users 120 always have enough data in the 15 buffer for transmit when selected by the allocation unit 114 and given a set of channelization codes for the transmission. To obtain a more general solution, additional constraints may be added on the set of assigned channelization codes, ϕ_m . These constraints include factors such as the data queue size of the user 120, and the target error rate requirements. Also, in the previous example, there are no specific restrictions imposed upon the assignment of the 20 channelization codes for each data user (i.e., other than they have to satisfy Equations (4), (5) and (13)).

Equation (13) can be rewritten as follows:

$$\sum_{m=1}^M \frac{\phi_m}{\gamma_{DPR} \hat{C}_{pilot,m}} = \frac{-N_c}{\lambda} - 1. \quad (17)$$

The optimization involves three linear equations with M variables. Therefore, for $M > 3$, there are more than one solution for the set $\{\phi_m\}$. By providing more structure in $\{\phi_m\}$ a scheduling algorithm for users with limited queue sizes and pre-defined target error rates may be determined.

One constraint imposed is to limit the number of modulation coding schemes (MCS) to a finite set, such as the set defined in Table 1, which shows an example of a finite set as a function of the information block size (denoted code block size), and the number of channelization codes. Table 1 is built on a multicode direct sequence CDMA (DS-CDMA) system, using a fixed spreading factor of 16. Exemplary parameters defined by the UMTS specification include a chip rate of 3.84Mcps and a slot duration of 0.67msec. Table 1 defines a pre-defined relation between the user buffer size, the number of channelization codes, and the MCSs that are possible under the given conditions. The allocation unit 114 takes these factors and relationships into account when assigning the available power and channelization codes to different users 120.

Number of Channel Codes	1280 bits code block	2560 bits code block	3840 bits code block	5120 bits code block	7680 bits code block	11520 bits code block	15360 bits code block
	640 Kbps	1280 Kbps	1920 Kbps	2560 Kbps	3840 Kbps	5760 Kbps	7680 Kbps
14	QPSK -6.98dB	QPSK -3.97dB	QPSK -2.48dB	QPSK -0.63dB	QPSK 1.29dB	16QAM 5.65dB	16QAM 8.1dB

12	QPSK -6.3dB	QPSK -3.34dB	QPSK -1.75dB	QPSK 0.34dB	QPSK 2.75dB	16QAM 6.85dB	16QAM 9.7dB
10	QPSK -5.4dB	QPSK -2.46dB	QPSK -0.73dB	QPSK 1.7dB	16QAM 5.18dB	16QAM 8.53dB	-
8	QPSK -4.5dB	QPSK -1.15dB	QPSK 0.80dB	QPSK 3.74dB	16QAM 6.86dB	-	-
6	QPSK -3.19dB	QPSK 0.55dB	QPSK 3.35dB	16QAM 6.02dB	16QAM 9.66dB	-	-
4	QPSK -1dB	QPSK 3.95dB	16QAM 7dB	-	-	-	-
2	QPSK 4.1dB	16QAM 9.9dB	-	-	-	-	-

Table 1: Modulation and Coding Schemes versus Channelization Codes, and the Required Es/Nt for 1% Packet Error Rate

After call admission control is performed at the core network (not shown), a set of users 120 are admitted with different QoS classes. For ease of illustration and to avoid obscuring the instant invention, it is assumed that any user 120 in the set has only one QoS value. The allocation unit 114 at the sector 130 then sub-divides the users 120 into different QoS categories. Again, users 120 from the same QoS class are scheduled and processed simultaneously, while users from different QoS classes are processed from highest quality down to the more delay tolerable classes separately.

In the following example, users 120 from the Same QoS class are scheduled. It is assumed that there are M data users in the same QoS class waiting to be scheduled and that there are no Hybrid-ARQ or IR operations initially. The allocation unit 114 first sorts the M users by their priorities in descending order according to some fairness criteria to form an ordered set, $U = \{u_i | u_1 \leq u_2 \leq \dots \leq u_M\}$. Exemplary fairness techniques known to those of ordinary skill in the art include round robin, maximum carrier to interference ratio (C/I), proportional fair, and the like. The allocation unit 114 then looks at the data queue 112 for the user 120 and determines the transmit code block sizes for each user 120. For example, if

user 1, u_1 , has 2000 bits in the data queue 112, then, the code block size of 1280 bits is selected for its next transmission. Based on the pre-determined target packet error rate, the allocation unit 114 then selects the MCS based on the estimates of the received signal-to-noise ratios (SNR) for each user 120. The estimate is computed using Equation (9) given the 5 pilot power and the code power for each channel are known at the sector 130. The received SNR estimates can be expressed as:

$$\hat{C}_m = \frac{|\eta_m|^2 \chi_m P_{data}}{I_{oc} + N_o}, \forall m. \quad (18)$$

Using Equation (18), the allocation unit determines the corresponding number of channelization codes to be used for the transmission by looking at the required SNR in the 10 finite set of MCSs for the selected code block size. For example, if user 1 has an estimated receive signal-to-noise ratio of -3dB , and a code block size of 1280 bits, four spreading codes will be assigned to user 1 for the next transmission, as shown in Table 1.

An iterative algorithm for solving the power and code assignment assuming that the user has been sorted by some fairness algorithm is now described with reference to the 15 simplified flow diagram of Figure 4. The code block sizes are determined for each user prior to entering the algorithm. In block 400, the set of data power fractions, $\{\chi_m^0\}$, is initialized to P_{data}/N_c for all m (*i.e.*, the available power is divided amongst the channelization codes). Of course other values for the initial power fraction assignments may be employed. In block 410, the value for the SNR estimates, \hat{C}_m , is determined in accordance with Equation (18). 20 In block 420, the values for the set of channelization codes for each user, $\{\phi_m^n\}$, $\forall \phi_m^n \in \mathbb{N}$, is determined from a finite rate table (*e.g.*, Table 1). In block 430, the value for λ is

determined in accordance with equation (13). In block 440, the values for the set of power fractions, $\{\chi_m^n\}$, is determined using Equation (11). In block 450, it is determined if the value of $\phi_m^n = \phi_m^{n+1}$, $\forall n > n_0$, where n_0 is sufficiently large (i.e., n_0 is a predetermined constant for the minimum number of iterations, e.g., 20, which may vary depending on the particular 5 implementation). If the solution has converged, the iteration terminates and the number of channelization codes is truncated at block 460. Otherwise, the iterative process repeats at block 410. Note that when selecting $\{\phi_m\}$, it is likely that the sum of all ϕ_m is larger than the maximum number of available channelization codes, N_c . Since all users 120 are sorted by priority, $\{\phi_m\}$ is truncated such that:

$$10 \quad \sum_{m=1}^{M_o} \phi_m + \phi'_{M_o+1} = N_c, \quad (19)$$

where $\phi'_{M_o+1} \leq \phi_{M_o+1}$ (i.e., set $\phi_m = 0$, $\forall m > M_o + 1$ and reduce the last non-zero terms in the set to meet the upper limit N_c). Also note that if there is reduction in ϕ_{M_o+1} then the transmission for the $(M_o + 1)$ user will likely have a higher packet error rate.

There are several different ways the allocation unit 114 may handle multiple QoS 15 classes, depending on whether the fairness algorithm in the allocation unit 114 is QoS aware and whether during call admission control (CAC) power constraints are specified for different QoS classes. Figures 5A through 5C show simplified block diagrams of three different techniques for scheduling users from multiple QoS classes.

The method illustrated in Figure 5A schedules users from the same QoS class first and 20 allocates power and channelization codes accordingly. In the method of Figure 5A, there are no predetermined proportions of the transmit power assigned to the various QoS classes. If

for the last user 120 the number of channelization codes or power is not sufficient to meet the target packet error rate, the unused power from the last user 120 in QoS class q may be used to schedule users 120 in QoS class $q+1$. Note that the allocation unit 114 may reserve extra power for the next QoS class, if desired, by taking power from the last few users 120 in the 5 current QoS class and using it for the next class.

In block 500, users 120 are admitted into QoS classes during CAC. In block 505, the users 120 in QoS class q are loaded into set Θ_q . In block 510, code block sizes are selected for packets in the set, Θ_q given P_{data} . In block 515, the users in Θ_q are sorted in accordance with a fairness algorithm. In block 520, resources are allocated in accordance with the 10 iterative method of Figure 5. In block 525, it is determined if the terminating condition, $\phi'_{M_0+1} = \phi_{M_0+1}$ is met (*i.e.*, all channelization codes and power are assigned). If the termination condition is met, the method ends in block 530. If resources still exist, the next QoS class is selected, $q = q + 1$, and the remaining power is determined (*i.e.*, $P_{data} = \chi_{M_0+1}^0 P_{data}$) in block 535. The scheduling operations then continue for the next QoS class at block 510.

15 The method illustrated in Figure 5B is similar to that of Figure 5A, but that in the method of Figure 5B, a pre-defined amount of power is assigned to each of the QoS classes. For example, during CAC portions of the total data power P_{data} may be assigned to the different QoS classes. Hence, after the users 120 are loaded into sets based on QoS class in block 505, the code block sizes are selected in block 540 for packets in the set, Θ_q given 20 $P_{data,q}$ (*i.e.*, the power reserved for the QoS class). Then, when the next QoS class is selected in block 545, the power need not be determined.

In the method illustrated in Figure 5C, the allocation unit is not constrained by QoS classes or power assignment for each QoS class. The allocation unit 114 processes users from all QoS classes at the same time, so there is no need to iterate over different QoS classes. However, the QoS class may be one of the factors employed in the fairness 5 algorithm to sort the users, so users in higher QoS may have higher relative positions in the sorted list. Users 120 are admitted into Q QoS classes in block 550. All the users 120 are sorted using a fairness algorithm in block 555. In block 560, the resources are allocated in accordance with the iterative method of Figure 5, and the method terminates in block 565.

For systems with Hybrid-ARQ and/or Incremental Redundancy, the allocation unit 10 may need to schedule users 120 with re-transmissions. During the re-transmission, the information size (*i.e.*, the code block size) for the packet to be re-transmitted is kept the same as previous transmission/re-transmissions and the allocation unit optimizes the channel capacity as described above.

The particular embodiments disclosed above are illustrative only, as the invention 15 may be modified and practiced in different but equivalent manners apparent to those skilled in the art having the benefit of the teachings herein. Furthermore, no limitations are intended to the details of construction or design herein shown, other than as described in the claims below. Consequently, the method, system and portions of the described method and system may be implemented in different locations, such as the wireless unit, the base station, a base 20 station controller and/or mobile switching center. Moreover, processing circuitry required to implement and use the described method and system may be implemented in application specific integrated circuits, software-driven processing circuitry, firmware, programmable logic devices, hardware, discrete components or arrangements of the above components as

would be understood by one of ordinary skill in the art with the benefit of this disclosure. It is therefore evident that the particular embodiments disclosed above may be altered or modified and all such variations are considered within the scope and spirit of the invention. Accordingly, the protection sought herein is as set forth in the claims below.